Enhancing FP-Growth Performance Using Multi-threading based on Comparative Study

Yousef K. Abu Samra, Ashraf Y. A. Maghari
Faculty of Information Technology
Islamic University of Gaza, Palestine

Abstract
The time required for generating frequent patterns plays an important role in mining association rules, especially when there exist a large number of patterns and/or long patterns. Association rule mining has been focused as a major challenge within the field of data mining in research for over a decade. Although tremendous progress has been made, algorithms still need improvements since databases are growing larger and larger. In this research we present a performance comparison between two frequent pattern extraction algorithms implemented in Java, they are the Recursive Elimination (RElim) and FP-Growth, these algorithms are used in finding frequent itemsets in the transaction database. We found that FP-growth outperformed RElim in term of execution time. In this context, multithreading is used to enhance the time efficiency of FP-growth algorithm. The results showed that multithreaded FP-growth is more efficient compared to single threaded FP-growth.

1. Introduction
The problem of mining frequent itemsets arose first as a sub-problem of mining association rules. While Apriori algorithm is quite successful for market based analysis in which transactions are large but frequent items generated is small in number [1]. Frequent itemsets play an essential role in many data mining tasks that try to find interesting patterns from databases such as association rules, correlations, sequences, classifiers, clusters and many more of which the mining of association rules is one of the most popular problems [2]. Also Sequential association rule mining is one of the possible methods to analysis of data used by frequent itemsets[3]. The original motivation for searching association rules came from the need to analyze so called supermarket transaction data, that is, to examine customer behavior in terms of the purchased products. Association rules describe how often items are purchased together. For example, an association rule “Bread, Cheese (60%)” states that four out of five customers that bought Bread also bought Cheese. Such rules can be useful for decisions concerning product pricing, promotions, store layout and many others.

Studies of Frequent Pattern Mining is acknowledged in the data mining field because of its importance. Hence, efficient algorithms for mining frequent patterns are crucial for mining association rules as well as for many other data mining tasks [1]. The major challenge found in frequent pattern mining is a large number of result patterns. As the minimum threshold becomes lower, an exponentially large number of patterns are generated. Therefore, pruning unimportant patterns can be done effectively in mining process and that becomes one of the main topics in frequent pattern mining. Consequently, the main aim is to optimize the process of finding patterns which should be efficient, scalable and can detect the important patterns which can be used in various ways.

The FP-growth algorithm transforms the problem of finding long frequent patterns to searching for shorter ones recursively and then concatenating the suffix [4]. It uses the least frequent items as a suffix, offering good selectivity. It is widely applied and performance studies demonstrate that the method substantially reduces search time.

There are many alternatives and extensions to the FP-growth approach, including depth-first generation of frequent itemsets [5], H-Mine [6], explores a hyper-structure mining of frequent patterns; building alternative trees; exploring top-down and bottom-up traversal of such trees in pattern-growth mining [7], and an array-based implementation of prefix-tree-structure for efficient pattern growth mining [8]. Recursive elimination [4] is one algorithm that does its work without prefix trees or any other complicated data structures, all the work is done in one simple recursive function.

DepthProject [5] mines only maximal frequent itemsets. It performs a mixed depth-first and breadth-first traversal of the itemset lattice. Also GenMax [9] is a backtrack search based algorithm for mining maximal frequent itemsets that uses a number of optimizations to prune the search space. H-Mine [6] is an algorithm for discovering frequent itemsets by a simple and novel data structure using hyper-links, H-struct, and a new mining algorithm, Hmine, which takes advantage of this data structure and dynamically adjusts links in the mining process.
Pincer-search algorithm [10], proposes a new approach for mining maximal frequent itemset which combines both bottom-up and top-down searches to identify frequent itemsets effectively. It classifies the data source into three classes as frequent, infrequent, and unclassified data.

FP-growth is one of the algorithms which is based upon the recursively divide and conquer strategy. FP-growth is proofed to be efficient and is widely used and applied. RElim, however, is proposed to be the algorithm of choice if implemented properly [4]. Moreover, RElim was inspired by H-Mine which is similar to FP-growth.

In this paper we compare FP-growth with RElim algorithm in term of time efficiency. We also use multithreading technique to enhance the time efficiency of FP-Growth algorithm. A comparison in terms of execution time were carried out between FP-growth and the enhanced FP-growth.

The rest of this paper is organized as follows: Section 2 discusses state of the art and reviews some related works. Section 3 explains the theoretical background needed. Section 4 describes our proposed multithreaded FP-growth. Section 5 shows experimental setup. Section 6 presents results and discussion. Finally Section 7 concludes the paper.

2. Related Work

Vina et al [1] provided a comparison between H-mine, Fp-Growth and SaM. A framework has developed to allow the flexible comparison of the algorithms. They measured time complexity and came to that the execution time of all the discussed algorithms is nearby but it can also be noticed that the execution time of SaM is comparatively less for higher support threshold.

Christian [4] presented a paper on Recursive Elimination algorithm. He proposed that if a quick and straightforward implementation is desired, it could be the method of choice. Even though its underlying scheme which is based on deleting items, recursive processing, and reassigning transactions is very simple and works without complicated data structures, recursive elimination performs surprisingly well.

Jochen Hipp et al. [11] provided several efficient algorithms that convoy with the popular and computationally expensive task of association rule mining with a comparison of these algorithms concerning efficiency. He proposed that the algorithms show quite similar runtime behavior in their experiments.

Aggarwal and Srikant [12] presented two new versions of Apriori, AprioriT and ApriorITID, for discovering all significant association rules between items in a large database of transactions and compared these algorithms to the previously known algorithms, the AIS and SETM algorithms. They proposed that these algorithms always outperform AIS and SETM.

Borgelt [13] provided efficient implementation of the more sophisticated approaches known under the names of Apriori and Eclat. Both rely on a top down search in the subset lattice of the items. He proposed for free item sets Eclat wins the competition with respect to execution time and it always wins with respect to memory usage. The data set in which it takes lead is for the lowest minimum support value tested, indicating that for lower minimum support values it is the method of choice, while for higher minimum support values its disadvantage is almost negligible. For closed item sets the more efficient filtering gives Apriori a clear edge with respect to execution time. For maximal item sets the picture is less clear. If the number of maximal item sets is high, Apriori wins due to its more efficient filtering, while Eclat wins for a lower number of maximal item sets due to its more efficient search.

Györi and Holban [14] had performed an experimental comparison between Apriori, DynFP-Growth, FP-growth, the algorithms were implemented in Java and tested on several data sets. They stated that FP-growth version out performed Apriori in all cases, and Apriori has the most memory consumption. On the other hand, the frequent database scans gave Apriori the maximum number of generated itemsets.

Shankar and Purusothaman [15] presented a comparative study of various methods in existence for frequent itemset mining, association rule mining with utility considerations. THUI (Temporal High Utility Itemsets)-Mine, heap mine (H-mine), and DSM-FI (Data Stream Mining for Frequent Itemsets) algorithms have been evaluated based on their memory usage for mining the frequent itemsets and association rules from large databases.

Vani [16] has conducted a Comparative Analysis of Association Rule Mining Algorithms Based on performance Survey between FP-growth and Eclat, as the fastest algorithms on the survey, and he concluded that their performance varies according to the data set used. In this paper, we compare a java implementation of Recursive Elimination and FP-Growth in term of execution time. We use two different datasets with different numbers of records and attributes, comparing their performance at low and high minimum supports.

High-performance parallel and distributed computing is becoming increasingly important as data keep growing in size and becoming complicated. Works have been done to parallelize the mining process such as shared memory systems. Rathi et al. [17] proposed a model that implements a parallel FP Growth algorithm that makes use of multiple Graphic processing (GPU) system, the proposed algorithm improves performance of the
algorithm. Wang and Wang [18] designed a parallel algorithm that works on distributed data framework, their algorithm does not need to create the whole FP tree, so it can handle huge data.

Frontier Expansion [19] is a new parallel Frequent Itemset Mining algorithm, its implementation can achieve good performance in heterogeneous platforms with shared memory multiprocessor and multiple Graphic Processing units and speedup 6-30 times sequential Eclat. Accelerating Parallel Frequent Itemset Mining on Graphic Processors with Sorting APFMS [20], is an algorithm that utilizes new generation GPUs to accelerate the mining process on opencL platform, results showed reduction in computation time.

Our work, however, is different in a way that it does not require any special prepared platforms nor hardware equipment such as GPUs. It can work on any computer that supports threading, and all computers do these days. By this we can accelerate the process of mining frequent patterns multiple times regardless of the hardware it runs on.

3. Theoretical Background

3.1. FP-growth

FP-tree algorithm is based upon the recursively divide and conquers strategy: first the set of frequent 1-itemset and their counts is discovered. With start from each frequent pattern, construct the conditional pattern base, then its conditional FP-tree is constructed (which is a prefix tree.). Until the resulting FP-tree is empty, or contains only one single path. (Single path will generate all the combinations of its sub-paths, each of which is a frequent pattern). The items in each transaction are processed in L order (i.e. items in the set were sorted based on their frequencies in the descending order to form a list) [3]. The detail step is as follows:

**FP-Growth Method:** Construction of FP-tree. Create root of the tree as a “null”. After scanning the database D for finding the 1-itemset and their counts is discovered. With start from each frequent pattern, construct the conditional pattern base, then its conditional FP-tree is constructed (which is a prefix tree.). Until the resulting FP-tree is empty, or contains only one single path. (Single path will generate all the combinations of its sub-paths, each of which is a frequent pattern). The items in each transaction are processed in L order (i.e. items in the set were sorted based on their frequencies in the descending order to form a list) [3]. The detail step is as follows:

**FP-Growth Method:** Construction of FP-tree. Create root of the tree as a “null”. After scanning the database D for finding the 1-itemset and their counts is discovered. With start from each frequent pattern, construct the conditional pattern base, then its conditional FP-tree is constructed (which is a prefix tree.). Until the resulting FP-tree is empty, or contains only one single path. (Single path will generate all the combinations of its sub-paths, each of which is a frequent pattern). The items in each transaction are processed in L order (i.e. items in the set were sorted based on their frequencies in the descending order to form a list) [3]. The detail step is as follows:

**Algorithm 1 (FP-tree construction)**

**Input:** A transaction database DB and a minimum support threshold.

**Output:** Its frequent pattern tree, FP-tree Method:

1. Scan the transaction database DB once. Collect the set of frequent items F and their supports. Sort F in support descending order as L, the list of frequent items.
2. Create a root of an FP-tree, T, and label it as ‘null’.
3. For each transaction Trans in DB do the following.
    1. Select and sort the frequent items in Trans according to the order of L. Let the sorted frequent item list in Trans be \([p \mid P]\), where \(p\) is the first element and \(P\) is the remaining list. Call insert tree\((p \mid P)\;T\).
   2. The function insert tree\((p \mid P)\;T\) is performed as follows.
      1. If \(T\) has a child \(N\) such that: \(N.item-name = p.item-name\), then increment \(N\)’s count by 1; else create a new node \(N\), and let its count be 1. Its parent link be linked to \(T\), and its node-link be linked to the nodes with the same itemname via the node-link structure. If \(P\) is nonempty, call insert tree\((P)\;N\) recursively.

**Algorithm 2 (FP-growth: Mining frequent patterns with FP-tree and by pattern fragment growth)**

**Input:** FP-tree constructed based on Algorithm 1, using DB and a minimum support threshold.

**Output:** The complete set of frequent patterns.

**Method:** Call FP-growth (FP-tree; null), which is implemented as follows.

**Procedure FP-growth (Tree; α)**

1. IF Tree contains a single path \(P\)
2. THEN FOR EACH combination (denoted as \(β\)) of the nodes in the path \(P\) DO
3. generate pattern \(βUα\) with support = minimum support of nodes in \(β\);
4. ELSE FOR EACH \(ai\) in header of Tree DO { (5) generate pattern \(β = aiUα\) with support = \(ai\).support; (6) Construct \(β\)’s conditional pattern base and then \(β\)’s conditional FP-tree Treeβ; (7) IF Treeβ ≠ (8) THEN Call FP-growth (Treeβ; β) }

3.2. Recursion elimination

In a pre-processing step delete all items from the transactions that are not frequent individually, i.e., do
not appear in a user-specified minimum number of transactions. This pre-processing is demonstrated in Figure 1, which shows an example transaction database on the left. The frequencies of the items in this database, sorted in an ascending order, are shown in the middle. If we are given a user-specified minimal support of 3 transactions, items f and g can be discarded. After doing so and sorting the items in each transaction in an ascending order by their frequencies we obtain the reduced database shown on the right of Figure 1.

Then select all transactions that contain the least frequent item (least frequent among those that are frequent), delete this item from them, and recurse to process the obtained reduced database, remembering that the item sets found in the recursion share the item as a prefix.

On return, remove the processed item also from the database of all transactions and start over, i.e., process the second frequent item etc.

![Figure 1. Transaction database (left), item frequencies (middle), and reduced transaction database with items in transactions sorted in an ascending order by their frequency (right) [21].](image)

This process is illustrated for the root level of the recursion, which shows the transaction list representation of the initial database at the very top, see Figure 2.

In the first step all item sets containing the item e are found by processing the leftmost list. The elements of this list are reassigned to the lists to the right (grey list elements) and copies are inserted into a second list array (shown on the right). This second list array is then processed recursively, before proceeding to the next list, i.e., the one for item a.

In these processing steps the prefix tree (or the H-struct), which is enhanced by links between the branches, is exploited to quickly find the transactions containing a given item and also to remove this item from the transactions after it has been processed.

![Figure 2. Procedure of the recursive elimination with the modification of the transaction lists (left) as well as the construction of the transaction lists for the recursion (right) [21].](image)

4. Proposed Multi-threaded FP-growth

Every path in the FP-tree keeps track of an itemset along with its support [4]. And it's known that each starting node generate it's related itemsets; for example consider the FP-growth tree in Figure 3., taking item 4 and 3 as an example it exists in three branches, below is a list of its related items along it's tree branch path.

1. Item 4
   - Branch1: {3,2,1}, supp 5
   - Branch2: {3,1}, supp 1
   - Branch3: {1}, supp 3

2. Item 3
   - Branch1: {2,1}, supp 5
   - Branch2: {1}, supp 2

So the candidate item sets of 4 are as follows:
- {3,2,1} Sup 5, {3,1} Sup 6, {1} Sup 9

and the candidate item sets of 3 are as follows:
- {2,1} Sup 5, {1} Sup 7

![Image of FP-growth tree](image)
The generated itemsets are concatenated with its support using (:) are

- **Itemsets of 4 are as follows:**
  - (4):9
  - (4,1):9
  - (4,2):5
  - (4,3):6
  - (4,1,2):5
  - (4,1,3):6
  - (4,2,3):5
  - (4,1,2,3):5

- **Itemsets of 3 are as follows:**
  - (3):7
  - (3,1):7
  - (3,2):5
  - (3,2,1):5

As we see from the previous example each item generates its own frequent itemsets, for example we can generate the itemset for item 4 and item 3 concurrently and then aggregate the results generated by both items in one file. So our contribution is to create a thread to gather the frequent itemsets of an item and after that combine the result of the finished thread in one single file.

The abstract code of our proposed threaded FP-growth can be found in the appendix.

5. **Experiment**

A. **System Information:**
   1) **Operating System:** windows server 2008, 64 bit
   2) **Memory:** 8192 MB Ram
   3) **Processor:** Intel® xeon® CPU ES-26200, 2.00GHz (4 CPUs).

B. **Implementation:**
   All the algorithms to be tested have been implemented in Java using eclipse 2015, with jdk8.0

C. **Datasets:**
   Two integer datasets were used in the comparison operation between ReLIm, FP-growth, and the proposed multithreaded FP-growth algorithms, we used integer data sets to be simple to deal with. These datasets contain different number of transactions and attributes.

   1) **Mushroom Dataset** [22]: it contains 8124 instances, this data set includes descriptions of hypothetical samples corresponding to 22 species of gilled mushrooms in the Agaricus and Lepiota Family.

   2) **Connect Dataset** [23]: it contains 67557 instance, 42 attributes. This database contains all legal 8-ply positions in the game of connect-4 in which neither player has won yet, and in which the next move is not forced.

D. **Optimization Issues:**
   Finding all frequent itemsets in a database is difficult since it involves searching all possible itemsets (item combinations). The set of possible itemsets is the power set over \( I \) and has size \( 2^n - 1 \) (excluding the empty set which is not a valid itemset).

   In this paper we assume that by using the multithreading technique we may enhance the execution time of finding the frequent item set in test datasets.

   Multithreading [24] is a programming technique for implementing application concurrency and, therefore, also a way to exploit the parallelism of shared memory multi-processors. A traditional “single threaded” process could be seen as a single flow of control (thread) associated one to one with a program counter, a stack to keep track of local variables, an address space and a set of resources. Multithreading programming allows one program to execute multiple tasks concurrently, by dividing it into multiple threads, so we will divide the operation on finding the frequent item sets to multiple threads, 4 threads precisely, one for each CPU, to benefit from the concurrency execution to minimize the time.

E. **Experiment:**
   1. The first experiment has been performed to compare performance of FP-growth with ReLIm using two different datasets.
   2. The second experiment is has been performed to see how much enhancement does multithreading add to FP-growth compared to single threaded version. Multithreaded FP-growth has been discussed in section V.

6. **Results and Discussion**

   1. Results of the first experiment show the superiority of FP-growth over ReLIm on mushroom dataset especially when the minimum support (mnsup) is low as shown in Figure 4. FP-growth outperformed ReLIm in all cases but the difference is dramatically reduced by the increase of the minimum support. We believe that depends on the nature of the dataset transactions, noticing that the number of frequent itemsets has reduced from hundreds of thousands to a few
thousands with a slight increase in the minsupp. This could be justified mentioning that RElim may perform better if a programming language that supports pointer handling is considered for implementation.

Also the superiority of FP-growth is clear in the larger dataset especially when the min support is low, see Figure 5.

These results are consistent with the suggestion of Cristian [21] that if a quick and straightforward implementation is applied it may improve the performance of the algorithm. We believe that the huge difference of execution time between the two algorithms refer to the mechanism of java code optimization, and especially the garbage collection mechanism.

Despite using GPUs as special hardware, they all parallelized the work as we did, but we used multithreading instead. And all the results showed improvements in the time execution of the algorithms.

**2. The second experiment proofs that the multithreaded FP-growth is much faster than FP-growth itself. Multithreaded FP-growth consumed around quarter the time consumed by FP-growth on both data sets, see Figure 6 and Figure 7. We believe that the result is proportional to the number of threads, as we used 4 threads, one for each CPU. These results came consistent to other findings in [17], [19], [20].**

**7. Conclusion and future work**

This paper presented a performance comparison between two frequent pattern extraction algorithms RElim and FP-Growth which implemented in Java. These two algorithms are used in finding frequent itemsets in the transaction database. We found out by experiment that FP-growth outperformed RElim in term of execution time. In this context, multithreading technique is used to enhance the time efficiency of FP-Growth algorithm in the part of generating frequent itemsets. The results showed that multithreaded FP-growth is more efficient compared to single thread FP-growth, this result comes clear at low minimum support thresholds.

In our experiment, it was difficult to accurately measure the space required by the algorithms because Java garbage collection mechanism make results about space unexplainable. So, in future work
we will try the two algorithms on another platform and compare the space complexity.

8. Appendix

The abstract code of our proposed FP-growth threading approach

```
/******************
-transactionCount is the total number of transactions
-mapSupport is hashmap of support of each item
-relativeMinsupp is the relative minimum support
-freqItemSet is ConcurrentHashMap which is a centralized container of the generated frequent itemsets
-item is the item which we need to generate its related itemsets.
-threads is a container of the running threads
******************/

int j=0; //counter of started threads
intThreadNum=4; // number of threads
FPGrowthTree tree = Constructs_FP_growth_tree();
for(inti=0;i<itemListSize;i++)
    if(j<ThreadNum)
        item = tree.headerList.get(i);
        CandidateGenerator c = new CandidateGenerator(tree,transactionCount,
            mapSupport, relativeMinsupp,
            freqItemSet,item);
        threads.add(c);
        c.start();
        j++;
    }
    if(j==ThreadNum || i==itemListSize){
        boolean exit=false;
        while(!exit){
            Iterator Itr = threads.iterator();
            while (Itr.hasNext()) {
                CandidateGenerator c = Itr.next();
                if(c.finished){
                    // a thread is finished, remove it from the threads list
                    Itr.remove();
                    //change exit flag to exit while loop
                    exit=true;
                }
            }
            if(exit){
                //if condition for check finished thread
                } // loop over the threads list
                } // loop to monitor the finished thread
                } // if condition to start threads monitoring
                } // main item loop
```

9. Acknowledgment

I would like to express my deep sense of gratitude and sincere thanks to my friend and colleague Ramzi A. Matar. His constructive and insightful comments, suggestions and help have improved the quality of this work.

10. References


